Computation Alignment:

Capacity Approximation without Noise Accumulation

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Capacity: $C = \log(1 + SNR)$

























K =Network Width

$D = \mathsf{Network} \mathsf{Depth}$

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 Compute-and-Forward (Nazer-Gastpar '11): Decode equations of messages. Partially captures signal interactions and removes noise.
What rates are attainable?

Compute-and-Forward to a Single Relay



- Decoder makes an estimate $\hat{\mathbf{u}}$ of a linear equation $\mathbf{u} = \bigoplus_{\ell=1}^{K} a_{\ell} \mathbf{w}_{\ell}$
- Require that probability of error vanishes with blocklength.
- Receiver can use its knowledge of the channel gains to match the integer equation coefficients a_{ℓ} to the channel coefficients h_{ℓ} .

Compute-and-Forward – Single Receiver

Lattice codes: Sums of codewords are codewords.

Erez-Zamir '04: Nested lattice codes can approach AWGN capacity.



Theorem (Nazer-Gastpar IT Trans. '11)

A relay can reliably decode the linear equation with coefficients $\mathbf{a} \in \mathbb{Z}^{K}$ from a channel with coefficients $\mathbf{h} \in \mathbb{R}^{K}$ if the message rates satisfy:

$$R_{\ell} < \max_{\alpha \in \mathbb{R}} \frac{1}{2} \log \left(\frac{\mathsf{SNR}}{|\alpha|^2 + \mathsf{SNR} \|\alpha \mathbf{h} - \mathbf{a}\|^2} \right)$$

For complex-valued channels, use real-valued decomposition.

Compute-and-Forward – Multiple Receivers



Theorem (Nazer-Gastpar IT Trans. '11)

The relays can decode a full rank set of linear equations from a channel matrix with rows $\mathbf{h}_k^T \in \mathbb{C}^K$ as long as for some full rank matrix with columns $\mathbf{a}_k \in \{\mathbb{Z} + j\mathbb{Z}\}^K$ the message rates satisfy:

$$R_{\ell} < \min_{k:a_{k\ell} \neq 0} \log \left(\frac{1 + \mathsf{SNR} \|\mathbf{h}_{k}\|^{2}}{\|\mathbf{a}_{k}\|^{2} + \mathsf{SNR}(\|\mathbf{h}_{k}\|^{2} \|\mathbf{a}_{k}\|^{2} - (\mathbf{h}_{k}^{*}\mathbf{a}_{k})^{2})} \right)$$

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Theorem (Niesen-Whiting ISIT '11, arXiv '11)

Coupling compute-and-forward with interference alignment can achieve K degrees-of-freedom.

• Key Idea: Use alignment scheme from Motahari et al. '09 to create effective integer channels that are good for compute-and-forward.

- In the high SNR limit, noise accumulation does not show up: we could just as well use compress-and-forward.
- Goal for this talk: Design a computation alignment scheme that can send equations with rates within a constant gap of $K \log SNR$.
- Ergodic interference alignment schemes offer excellent finite SNR rates for the standard interference channel problem.
- Requires time-varying channel coefficients. Our constant gap will depend on the number of sources K and the fading statistics but not the network depth D.

Computation Alignment - Setup



• Relays recover an invertible set of functions

$$u_k = f_k(w_1, w_2, \dots, w_K)$$

- Vanishing probability of error.
- Channel coefficients are independent from each other, stationary and ergodic across time, and have uniform phase.

Basic ergodic alignment scheme:

- Assume each relay wants a single message, $u_k = w_k$.
- Nazer-Gastpar-Jafar-Vishwanath '09: The following rate is achievable for each user:

$$R_k = \frac{1}{2} \mathbb{E} \left[\log(1 + 2|h_{kk}|^2 \mathsf{SNR}) \right]$$

• Corresponds to a sum rate that scales like $\frac{K}{2}\log \text{SNR}$.

1. At time t with channel **H**, user k transmits symbol X_k .

$$\mathbf{H} = \begin{bmatrix} h_{11} & h_{12} & \cdots & h_{1K} \\ h_{21} & h_{22} & \cdots & h_{2K} \\ \vdots & \vdots & \ddots & \vdots \\ h_{K1} & h_{K2} & \cdots & h_{KK} \end{bmatrix}$$

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2. When complementary matrix \mathbf{H}_C occurs, retransmit symbol X_k .

$$\mathbf{H}_{C} = \begin{bmatrix} h_{11} & -h_{12} & \cdots & -h_{1K} \\ -h_{21} & h_{22} & \cdots & -h_{2K} \\ \vdots & \vdots & \ddots & \vdots \\ -h_{K1} & -h_{K2} & \cdots & h_{KK} \end{bmatrix}$$

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3. Otherwise, transmit new symbols and wait for their \mathbf{H}_C .

- For long block lengths, nearly every matrix finds its quantized match.
- In general, we can successfully pair together matrices if their coefficients are phase rotations of each other,

$$h_{k\ell}[t_C] = e^{j\phi_{k\ell}} h_{k\ell}[t] \; .$$

• Computation alignment: Pair up coefficients to create effective integer channels.

Computation Alignment - Motivating Example



• Send 3 symbols over 2 channel uses. Sum rate is:

$$R_{\mathsf{sum}} = \frac{3}{2} \mathbb{E} \left[\min_{k,\ell} \log(1 + |h_{k\ell}|^2 \mathsf{SNR}) \right]$$

- Each relay gets 2 equations. Coefficients *a* and *b* are non-zero integers so equations are always invertible.
- Use compute-and-forward to code over resulting effective channels.









$$Y_1[t] - Y_1[t_C] =$$







$$X_1[t] = S_{1A} + S_{1B}$$
$$X_1[t_C] = S_{1A} - S_{1B}$$

















 $Y_1[t] + Y_1[t_C] =$ $Y_1[t] - Y_1[t_C] =$







 $Y_1[t] + Y_1[t_C] = 2h_{11}S_{1A}$ $Y_1[t] - Y_1[t_C] = 2h_{11}S_{1B}$

















$$\begin{split} Y_1[t] + Y_1[t_C] &= 2(h_{11}S_{1A} + h_{12}S_2) \\ Y_1[t] - Y_1[t_C] &= 2h_{11}S_{1B} \end{split}$$



Adjust strength of S_{1A} to get integer coefficients.

$$X_1[t] = S_{1A} + S_{1B}$$
$$X_1[t_C] = S_{1A} - S_{1B}$$

Tx 1

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Tx 1 Ad to ge

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Adjust strength of S_{1B} to get integer coefficients.

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Rx 1





$$\begin{split} Y_2[t] + Y_2[t_C] &= 2(h_{21}S_{1B} + h_{22}S_2) \\ Y_2[t] - Y_2[t_C] &= 2h_{21}\rho S_{1A} \end{split}$$

Adjust strength of S_{1B} to get integer coefficients.

$$X_1[t] = \rho S_{1A} + \gamma S_{1B}$$
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Tx 1

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Rx 1





$$\begin{split} Y_2[t] + Y_2[t_C] &= 2h_{22}(bS_{1B} + S_2) \\ Y_2[t] - Y_2[t_C] &= 2h_{21}\rho S_{1A} \end{split}$$

Computation Alignment - Main Result

K = 2 case:

- User 1 sends L symbols using L DFT vectors.
- User 2 sends L-1 symbols using L-1 DFT vectors.
- Relays get 2L equations. Sum rate is approximately $\frac{2L-1}{L} \log SNR$.

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 $K \ge 2$ case: More sophisticated set of alignment vectors.

Theorem (Niesen-Nazer-Whiting arXiv '11)

For Rayleigh fading, $h_{k\ell}[t] \sim C\mathcal{N}(0,1)$, the following sum rate is achievable using computation alignment:

$$R_{sum} = K \log \mathsf{SNR} - 7K^3$$
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- Penalty term is due to setting rates using the weakest coefficient.
- For equal magnitudes, $|h_{k\ell}| = 1$, we can achieve $R_{sum} = K \log SNR$.



Avestimehr-Diggavi-Tse '11: $C \ge K \log SNR - g_1(K, D)$

This talk: $C \ge K \log SNR - g_2(K, fading)$



Avestimehr-Diggavi-Tse '11: $C \ge K \log SNR - KD$

This talk: $C \ge K \log SNR - 7K^3$

- Compute-and-forward combined with alignment as a new tool for capacity approximation.
- More work is needed to get a constant gap that does not depend on the fading statistics.
- Multiple receivers. Layered interference channels.
- Fixed channel coefficients.
- For more info: ISIT '11 tutorial slides on Algebraic Structure in Network Information Theory available at iss.bu.edu/bobak